Parallel Performance of Structured AMR Calculations Using the SAMRAI Framework

Andrew Wissink

with

Richard Hornung, Scott Kohn, David Hysom, Steve Smith, Noah Elliott, Brian Gunney

Center for Applied Scientific Computing Lawrence Livermore National Laboratory

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Outline

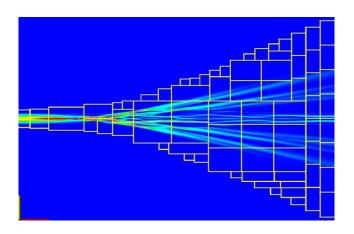
- SAMRAI introduction
- Parallel implementation of SAMR
- Parallel performance measurements
- New algorithms to enhance parallel performance
- Requirements and issues on BG/L

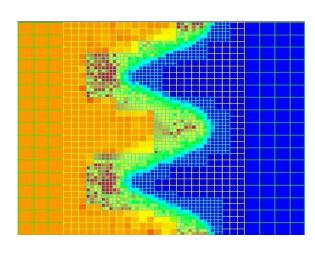
SAMRAI

Structured Adaptive Mesh Refinement Application Infrastructure

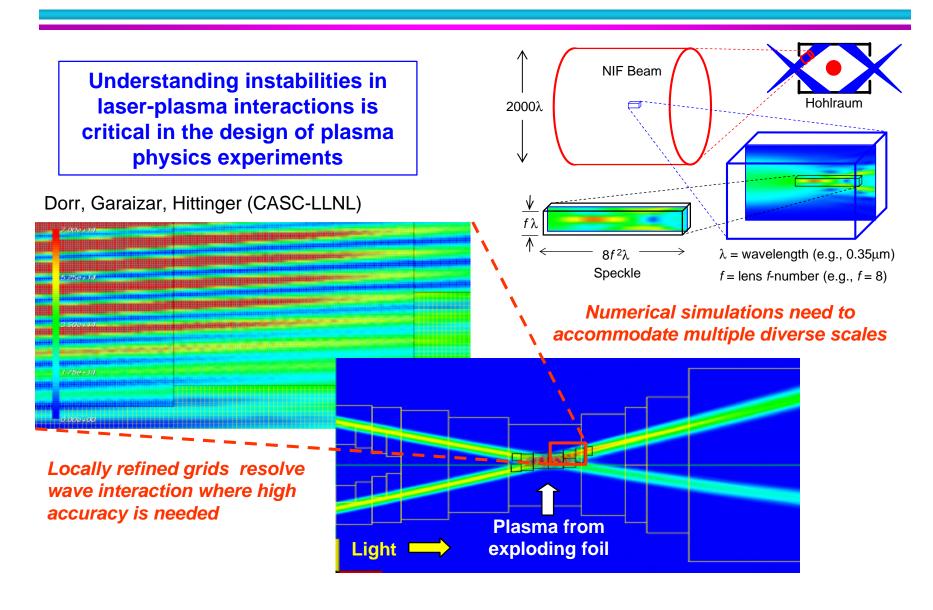
- Object-oriented (C++) software framework for parallel (MPI) adaptive multi-physics applications
- Supports applications investigating multi-scale phenomena.
- High-level reusable code and algorithms shared across a variety of applications.

www.llnl.gov/CASC/SAMRAI





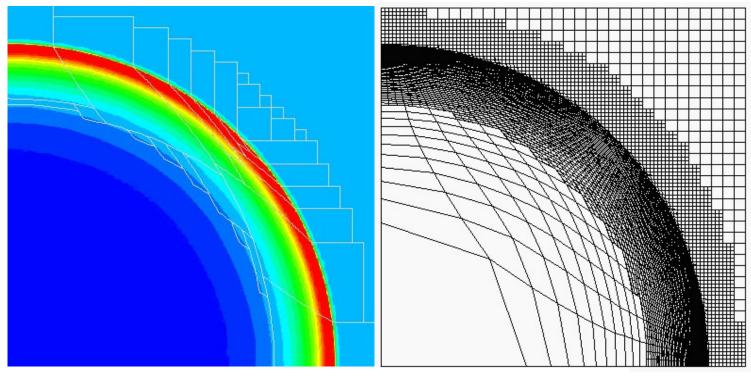
ALPS uses SAMRAI for adaptive laser plasma instability simulation



ALE-AMR couples ALE models with AMR to model shock hydrodynamics

Improve accuracy of ALE simulations by increasing concentration of mesh points around regions of interest

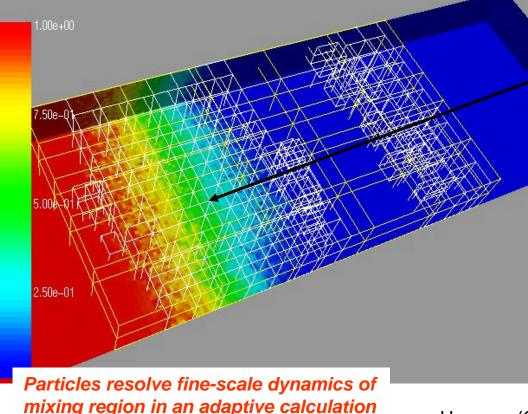
Anderson, Pember, Elliott (CASC-LLNL)



Sedov blast wave density and Lagrangian mesh

Hybrid continuum-DSMC model used to efficiently resolve interface dynamics

Interface instability problems (e.g., Richtmyer-Meshkov) involve coarse-scale hydrodynamic transport and fine-scale molecular diffusion



fluid A fluid B

Continuum representation
(Euler, Navier-Stokes)
away from interface

DSMC representation at interface

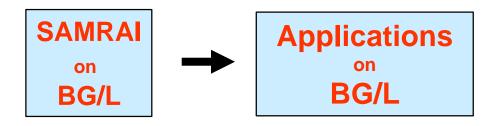
- Interface region grows and moves as instability evolves
- Standard CFD simulation of turbulent mixing is limited by finest mesh scale
- Particle resolve molecular behavior but are too expensive for large domains

Hornung (CASC), Garcia (SJSU)

shock

SAMRAI provides infrastructure support for a variety of applications

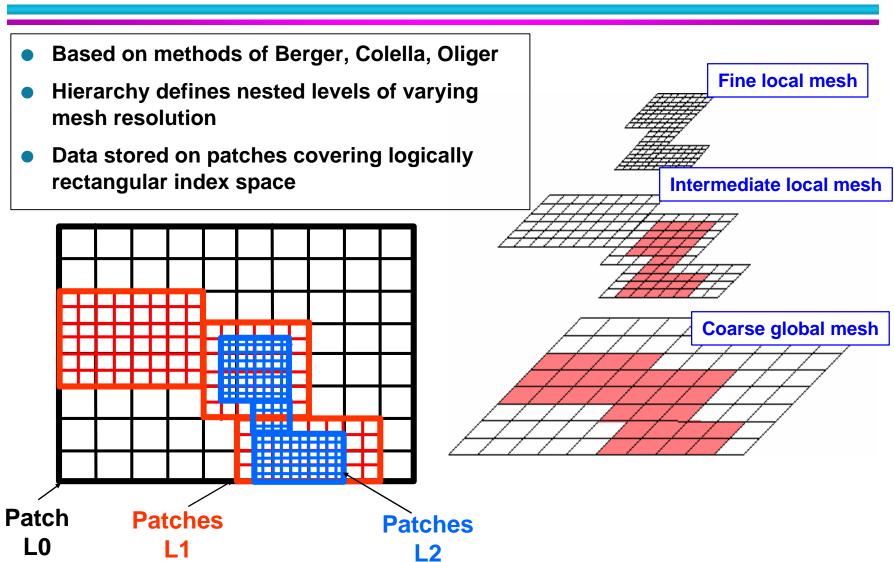
- Parallel processing support (MPI)
- Shared algorithms
- Interfaces for SAMR data to solvers (PETSc, PVODE, hypre)
- Checkpointing & restart support (HDF)
- Parallel tools (VAMPIR, TAU)
- Current users regularly run on existing large processor systems



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Structured AMR (SAMR) employs a dynamically adaptive "patch" hierarchy



Dynamic mesh adapts to features as solution evolves

Adaptive solution of **Euler equations**

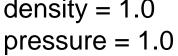
Initial conditions:

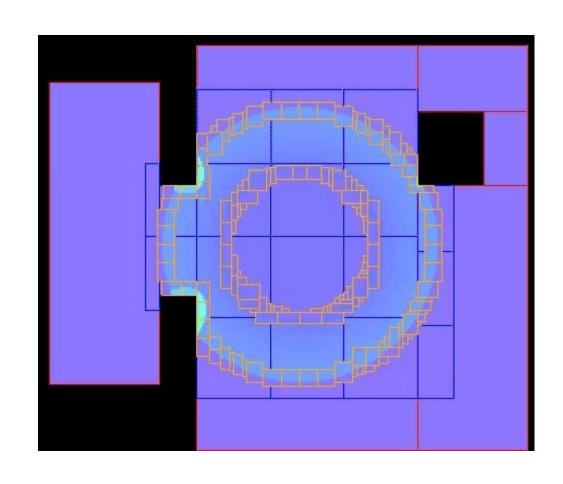
inside sphere

density = 8.0pressure = 40.0

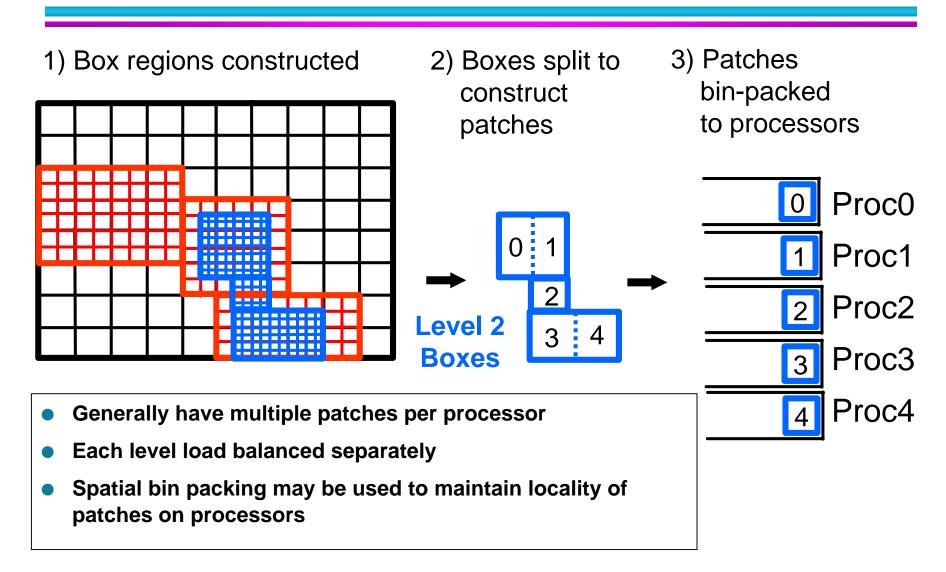
outside sphere

density = 1.0



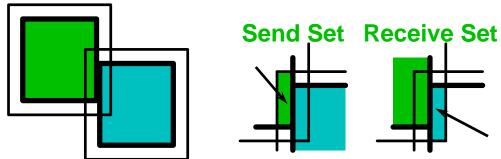


Patches distributed to processors to balance computational workload

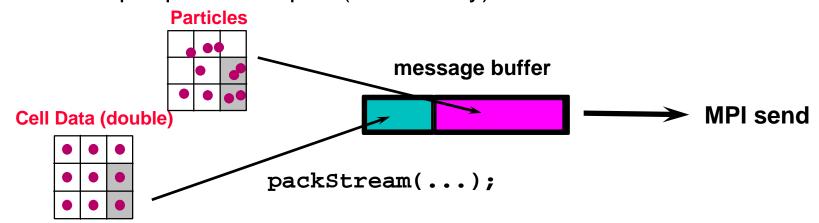


Communication schedules create and store data dependencies

 Amortize cost of creating send/receive sets over multiple communication cycles



- Data from various sources packed into single message stream
 - supports complicated variable-length data
 - one send per processor pair (low latency)



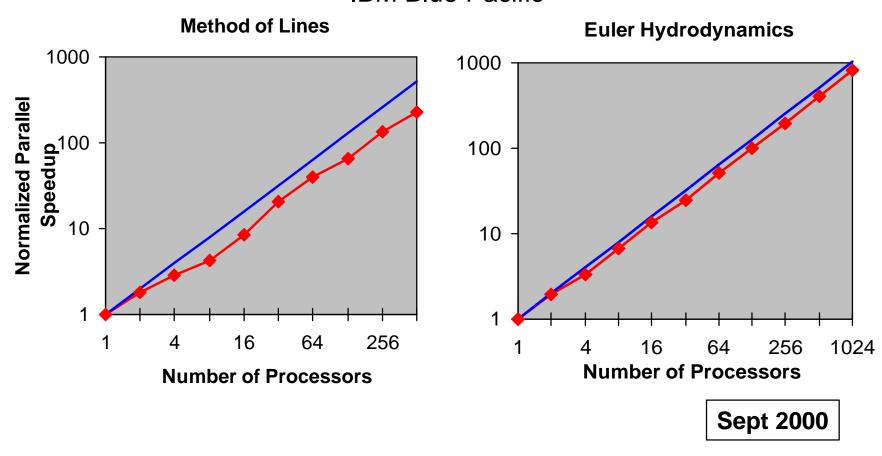
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Non-adaptive calculations using SAMRAI show good scaling

Single Level calculation

50x50x50 patch per processor IBM Blue Pacific



Benchmarks constructed to analyze scaling properties of SAMR applications

- Simple numerical kernels
- Invoke the main algorithmic components used in more complex apps (e.g. meshing, time advance, etc.)
- Timing decomposed into two phases:

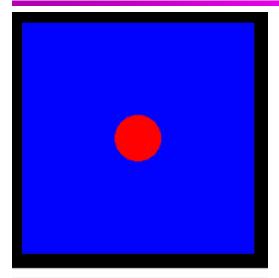
Time Advance:

- numerical kernels
- communication (filling ghost cells)
- load imbalances

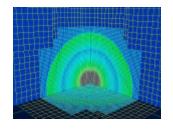
Re-Gridding:

- cluster tagged cells
- construct communication schedules
- distributing data to new mesh configuration

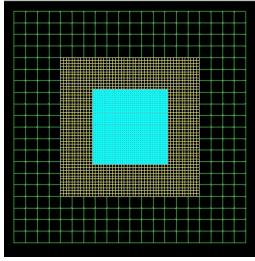
Non-scaled Euler benchmark – same problem size run on all processors

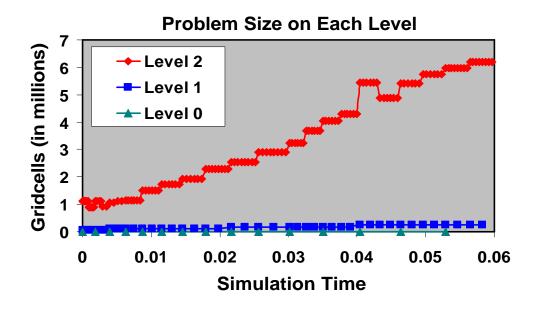


3D spherical shock - Euler hydrodynamics

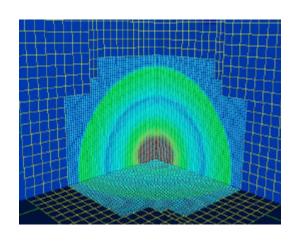


- Workload changes over simulation
- Per-processor workload decreases as number of processors increased



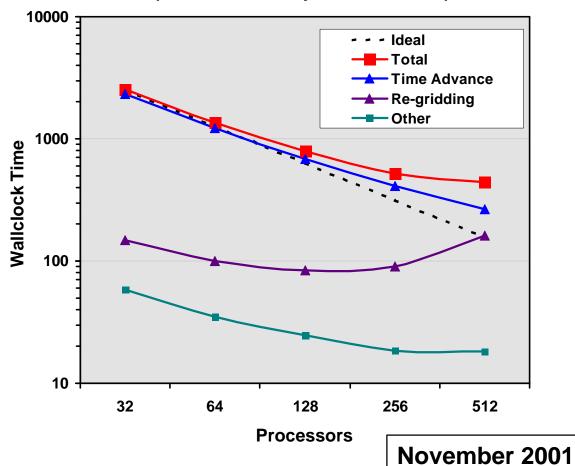


Parallel Performance of *non-scaled* adaptive Euler benchmark

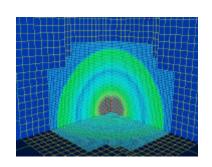


Non-scaled
Euler calculation
ASCI IBM Blue Pacific

Measured Solution Time on Various Processors (3 Level Euler Sphere Problem)



Poor scaling in re-gridding hurts efficiency on large processor counts

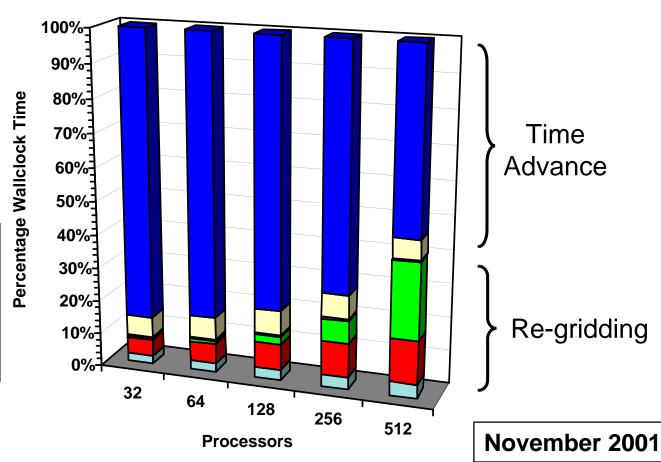


Non-scaled Euler calculation

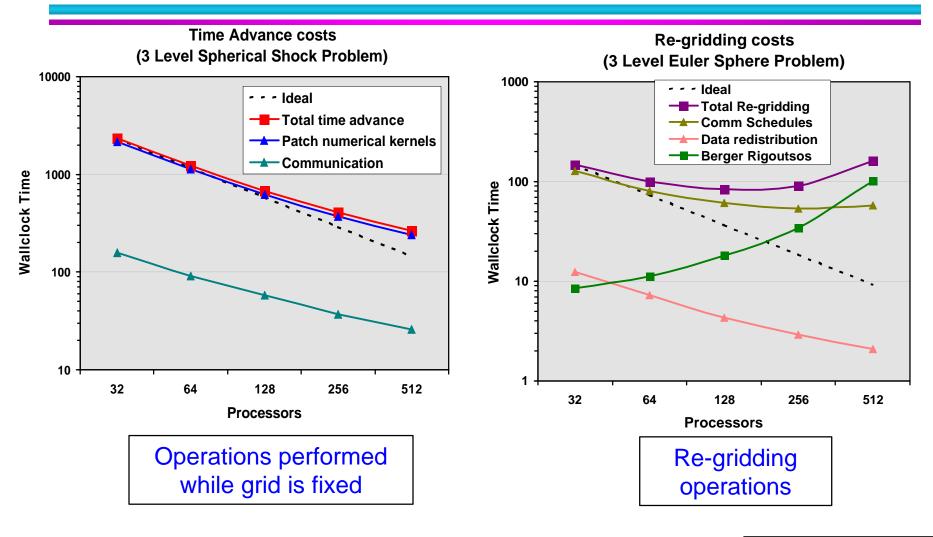
ASCI IBM Blue Pacific



Measured Solution Time on Various Processors (3 Level Spherical Shock Problem)

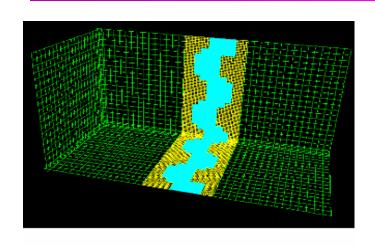


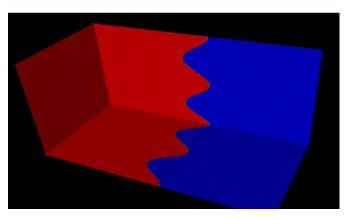
Poor scaling in re-gridding hurts efficiency on large processor counts (ASCI Blue Pac)



November 2001

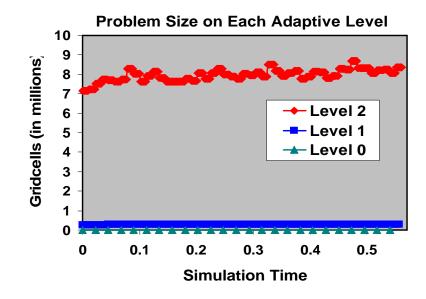
Scaled linear advection benchmark – problem size increased with processors



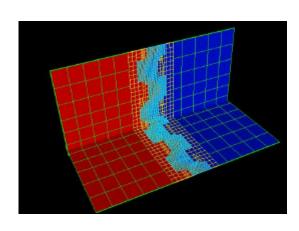


3D advecting sinusoidal front - linear advection

- Workload uniform over simulation
- Per-processor workload remains constant as number of processors is increased

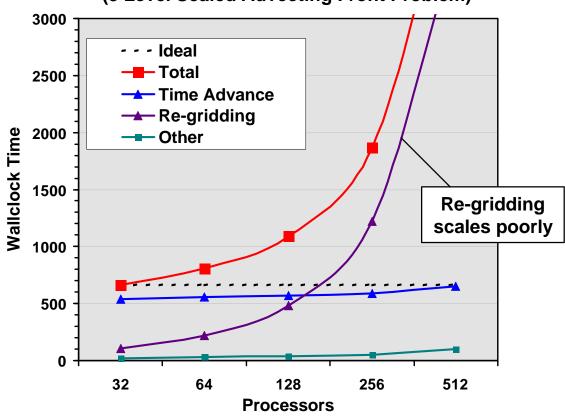


Parallel performance of scaled linear advection benchmark



Scaled
Linear advection
calculation
ASCI IBM Blue Pacific

Measured Solution Time on Various Processors (3 Level Scaled Advecting Front Problem)



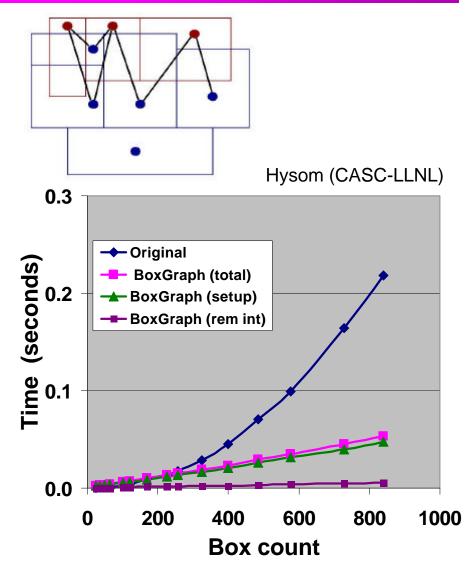
November 2001

Outline

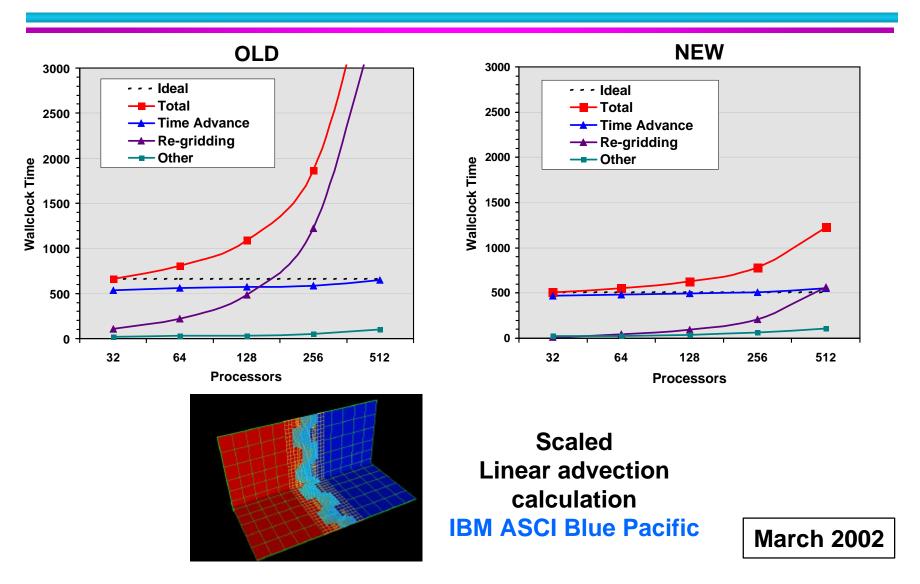
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Graph-based algorithms speed up communication schedule construction

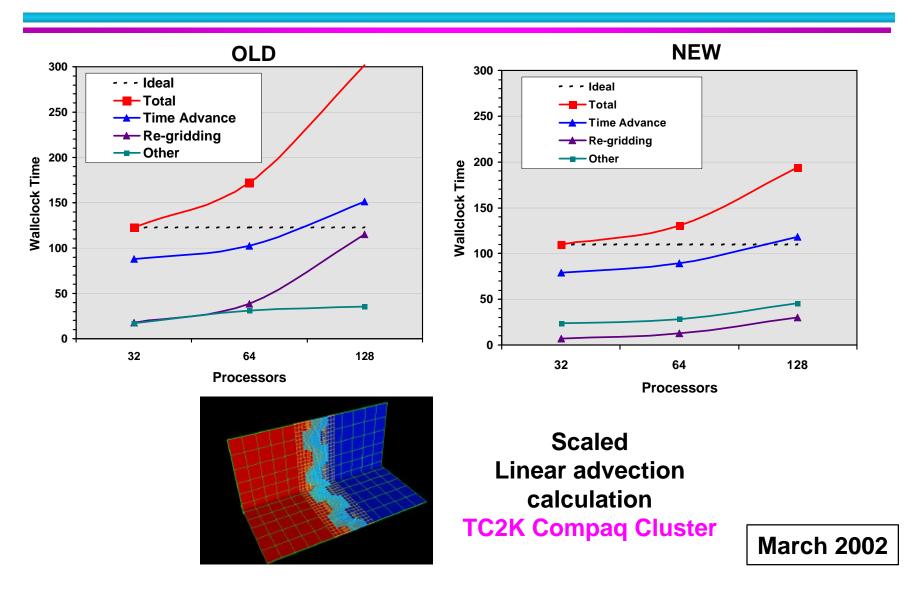
- Before constructing a communication schedule to transfer data between two levels, build a "Box graph":
 - Insert a vertex in V for each box
 - Insert an edge (i,j) at intersection
- Given this graph, can find a box's neighbors in O(1)
- Primary cost is graph construction



Scaled results with new graph-based schedule construction algorithm



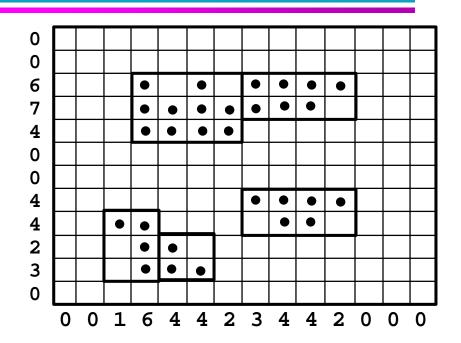
Scaled results with new graph-based schedule construction algorithm



Binary Tree reduction for tagged-cell clustering algorithm (Berger-Rigoutsos)

Berger-Rigoutsos:

- Forms new patches from tagged cells
- Determines box cuts from global histogram through recursive algorithm
- Original implementation used global reductions to form histogram

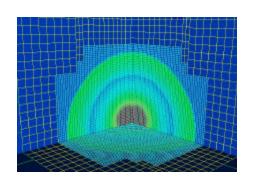


New Implementation:

- Binary-tree reduction algorithm collects information from selected processors at each recursion.
- New box configuration constructed and broadcast by one processor.

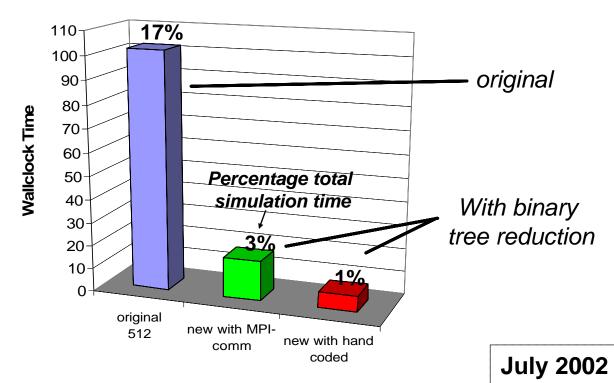
Timing results Berger-Rigoutsos algorithm with binary-tree reductions.

- Binary tree reduction algorithm two implementations
 - MPI communicators
 - Hand-coded MPI send/recvs



Non-Scaled
Euler calculation
IBM ASCI Blue Pacific

Berger-Rigoutsos – 512 processors



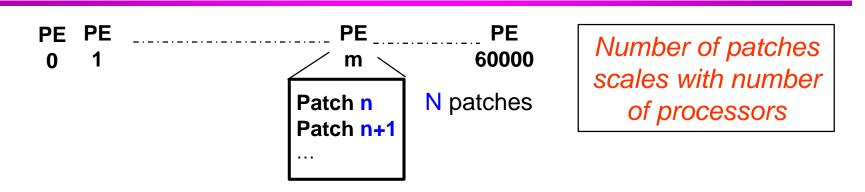
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BlueGene/L wish list

- SAMRAI Dependencies:
 - C++, C, F77/F90 compilers
 - MPI
 - HDF5 (checkpointing)
- Desirable features:
 - C++-capable debugging tool
 - Memory analysis tool (i.e. reports stack/heap usage on nodes)

Scaling issues with a large number of processors



- Box operations in gridding may invoke O(N²) algorithms (e.g. former communication schedule algorithm).
- More efficient graph-based algorithms work on up to 512 processors, but need to develop efficient algorithms for O(10K) processors.
- Difficult to assess beforehand because complexity is generally problem dependent.

BG/L will require us to rethink our grid storage approach

 Current approach: Each MPI process holds a box for each patch in the problem to determine communication dependencies.

 Because # patches grows with # processors, trivial overhead becomes non-trivial on BG/L.

		per processor	
procs	<u>patches</u>	storage (MB)	Large overhead for nodes of BG/L
0.5K 60K	2.5K-10K 300K-1200K	< 1 MB 20-80MB	

Collective communications on BG/L

- Berger-Rigoutsos clustering:
 - Binary tree reduction algorithm effective in reducing costs on O(0.5K) processors.
 - Will this approach be effective on O(10K) processors?
- Some global communications are necessary (e.g. timestep synchronization in time advance).

Concluding Remarks

- Porting SAMRAI to BG/L enables a variety of applications to use the architecture.
- Results of scaling AMR algorithms on up to 512 processors:
 - Communication not the primary source of scaling inefficiency.
 - Re-gridding operations that are trivial on small numbers of processors become significant on large numbers.
 - More efficient graph-based algorithms successful in reducing these costs.
- Speculation on running AMR applications on BG/L:
 - Re-gridding costs will likely be the main hindrance.
 - Continued exploration into more efficient gridding algorithms needed.

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- Document UCRL-PRES-149437